Computer-Aided Generation of Enforcement Mechanisms for Error-Tolerant Policies

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This talk is about…

runtime monitor

Policy

Application
In e-health scenario…

- **Users:**
  - doctors, nurses, pharmacists

- **Application:**
  - implementing a business process

- **Policy:**
  - default protocol
    - simple
    - well understood by the users
  - health and privacy regulations

- **Runtime monitor enforcing policy**
Once the runtime monitor is deployed...

- it ensures compliance
- when an insignificant deviation from the protocol happens
  - runtime exception etc. disturbs the users
  - they start being convinced that “The system doesn’t work”
  - if such disruptions occur too often, then...
  - the user will try to bypass the runtime monitor
- => some of the policy violations might be amended
How to specify such deviations?

- With all possible deviations in the policy?
- No. It will make policy too complicated for the end users!
- The policy should remain simple and well-understood
- Alternative: let runtime monitor
  - enforce simple policies
  - tolerate insignificant deviations
Case study: drug dispensation process from Hospital San Raffaele (Milan)

- high-level business process
- must comply with specific rules (File F), e.g.
  - if the patient is involved in the clinical trial for testing the drug,
    - the doctor should insert a special research protocol number
    - in order to reimburse this drug later by the clinical trial funds.
  - if the drug is highly sensitive,
    - the doctor should review the therapeutical notes
Main steps of the default protocol

- The doctor selects a drug for his patient.

- If the drug is highly sensitive, the doctor reviews therapeutical notes. Otherwise therapeutical notes can be emitted.

- In case the drug is registered for research program, the doctor should insert the research protocol number. Otherwise the doctor skips this step.

- The doctor performs other actions needed for the drug prescription.
Venial errors

- The doctor forgot to review the therapeutical notes
- (he closes the window instead of clicking “Done”)
- This error is called a **venial error**
- But the doctor should not be allowed to violate the policy systematically
- it should be allowed a limited number of times, for example $k$ times per day
The doctor did not insert the research protocol number for the research drug

Then the whole reimbursement process for this drug can go wrong!

This should be “corrected” by inserting another special number that will be used during audit

This error is called an amendable error

This number of corrections should be limited as well
More formally…

- Deviations are given with expected $ex$ and correcting $c$ functions

<table>
<thead>
<tr>
<th>Deviation $e$</th>
<th>Expected action $ex(e)$</th>
<th>Correction $c(e)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>The doctor did not review the therapeutical notes</td>
<td>The doctor reviews therapeutical notes</td>
<td>The doctor did not review the therapeutical notes</td>
</tr>
<tr>
<td>The doctor did not insert the research protocol number for the research drug</td>
<td>The doctor inserts the research protocol number</td>
<td>Insert a special number for auditing</td>
</tr>
</tbody>
</table>
More formally…

- Default protocol represented by a finite state automaton (policy automaton)

- Number of possible errors $k$
Runtime monitor model

- Execution transformer:
  - edit automaton (EA) (+ memory)
  - can insert and suppress actions

![Diagram showing the execution transformer model]
Automatic construction of edit automaton

- **Automatic construction:**
  - **Input:** default protocol, number of errors $k$, functions for deviations $ex$ and $c$,
  - **Output:** edit automaton

- State is augmented with an error counter $s$
- if an error $e$ is venial ($c(e) = e$), accept $e$ and increase $s$
- if an error $e$ is amendable ($c(e) \neq e$), accept $c(e)$ and increase $s$
- when an accepting state of default protocol is reached, reset the counter $s$
### Automatic construction of edit automaton

<table>
<thead>
<tr>
<th>s</th>
<th>user action e</th>
<th>is allowed?</th>
<th>s&lt;k</th>
<th>c(e) =e?</th>
<th>result</th>
<th>s’</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>The doctor selects the drug</td>
<td>yes</td>
<td></td>
<td></td>
<td>proceed to next state with e</td>
<td>0</td>
</tr>
<tr>
<td>0</td>
<td>The doctor did not review the therapeutical notes</td>
<td>no</td>
<td>yes</td>
<td>yes</td>
<td>proceed to next state with e</td>
<td>1</td>
</tr>
<tr>
<td>1</td>
<td>The doctor did not insert the research protocol number for the research drug</td>
<td>no</td>
<td>yes</td>
<td>no</td>
<td>proceed to next state with c(e)</td>
<td>2</td>
</tr>
<tr>
<td>2</td>
<td>another error</td>
<td>no</td>
<td>no</td>
<td>yes/no</td>
<td>move to error state and reset the counter</td>
<td>0</td>
</tr>
</tbody>
</table>
Formal properties: soundness & transparency

- Assume:
  - $\sigma$ is an actual execution
  - $P(\sigma)$ is true if an actual execution complies with the default protocol
  - $EA(\sigma)=\sigma'$ edit automaton transforms an actual execution
Soundness & transparency

- An enforcement mechanism EA is **sound** with respect to $P$ if for all $\sigma$: $P(EA(\sigma))$
- An enforcement mechanism EA is **transparent** with respect to $P$ if for all $\sigma$: $P(\sigma) \Rightarrow EA(\sigma)=\sigma$

actual executions $\sigma$  

$EA(\sigma)$
Theorem 1 [Transparency]
Given $<P, k, c, \text{ex}>$, the constructed enforcement mechanism EA is \text{transparent} with respect to $P$. 
Formal properties: predictability

- Predictable within $\varepsilon$
  - if the execution $\sigma$ is close to a valid execution $\sigma'$, then $EA(\sigma)$ should be close to $\sigma'$
Predictable within $\varepsilon$

- if the execution $\sigma$ is close to a valid execution $\sigma'$, then $EA(\sigma)$ should be close to $\sigma'$

actual executions $\sigma$

$E A(\sigma)$
Formal properties: predictability

- Predictable within $\varepsilon$
  - if the execution $\sigma$ is close to a valid execution $\sigma'$, then $EA(\sigma)$ should be close to $\sigma'$

- Metrics to define “close” relation:
  - replacing distance $d$ between executions

- Definition:
  - An enforcement mechanism $EA$ is **predictable within** $\varepsilon$ with respect to $P$, if for all $\sigma'$ such that $P(\sigma')$:
    
    for all $\nu \geq \varepsilon$ exists $\delta > 0$ such that for all $\sigma$:
    
    $d(\sigma, \sigma') \leq \delta \implies d(EA(\sigma), EA(\sigma')) \leq \nu$
Theorem 2 [Predictability]
Given \( \langle P, k, c, ex \rangle \), a constructed enforcement mechanism EA is \textit{predictable within } k.
Conclusions

- Policies must stay simple
- Some deviations should be tolerated
- An enforcement mechanism modeled as edit automata automatically constructed for it
- Proven to be transparent and k-predictable (for tolerating k errors)